

LAWVERE COMPLETION AND SEPARATION VIA CLOSURE

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Dedicated to Bill Lawvere at the occasion of his seventieth birthday

ABSTRACT. For a quantale V , first a closure-theoretic approach to completeness and separation in V -categories is presented. This approach is then generalized to \mathcal{T} -categories, where \mathcal{T} is a topological theory that entails a set monad \mathbb{T} and a compatible \mathbb{T} -algebra structure on V .

INTRODUCTION

Bill Lawvere’s 1973 milestone paper “Metric spaces, generalized logic, and closed categories” helped us to detect categorical structures in previously unexpected surroundings. His revolutionary idea was not only to regard individual metric spaces as categories (enriched over the monoidal-closed category given by the non-negative extended real half-line, with arrows provided by \geq and tensor by $+$), but also to expose the purely categorical nature of the key concept of the theory, Cauchy completeness. The first step to this end was to disregard metric conditions that actually obscure the categorical intuition. In fact, once one has dropped the symmetry requirement it seems much more natural to regard the metric d of a space X as the categorical hom and, given a Cauchy sequence (a_n) in X , to associate with it the pair of functions

$$\varphi(x) = \lim d(a_n, x) \quad \text{and} \quad \psi(x) = \lim d(x, a_n).$$

Lawvere’s great insight was to expose these functions as pairs of adjoint (bi)modules whose representability as

$$\varphi(x) = d(a, x) \quad \text{and} \quad \psi(x) = d(x, a)$$

is facilitated precisely by a limit a for (a_n) . Hence, a new notion of completeness for categories enriched over any symmetric monoidal-closed category V was born. Also in the enriched category context it is often referred to as Cauchy completeness. But since Lawvere’s brilliant notion entails no sequences at all, just the representability requirement for bimodules, this name seems to be far-fetched and, contrary to popular belief, was in fact not proposed in his paper. Hence, here we use *L-completeness* instead.

In the first part of this paper we give a quick introduction to V -category theory (see [Kel82]) in the special case of a commutative unital quantale V , focussing on the themes of *L-completion* and *L-separation*. We are not aware of an explicit prior occurrence of the latter notion, and both themes are treated with the help of a new closure operator that arises most naturally in the 2-category $V\text{-Cat}$, as follows. Call a V -functor $m : M \rightarrow X$ *L-dense* if $f \cdot m = g \cdot m$ implies $f \cong g$ for all V -functors $f, g : X \rightarrow Y$; the *L-closure* of a subobject M of X is then the largest subobject \overline{M} of X for which $M \rightarrow \overline{M}$ is *L-dense*. For *L-separated* V -categories, *L-dense* simply means epimorphism. The *L-separated reflection* of

2000 *Mathematics Subject Classification.* 18A05, 18D15, 18D20, 18B35, 18C15, 54B30, 54A20.

Key words and phrases. Quantale, V -category, monad, topological theory, module, Yoneda lemma, closure operator, completeness.

The first author acknowledges partial financial assistance by Unidade de Investigação e Desenvolvimento Matemática e Aplicações da Universidade de Aveiro/FCT.

The second author acknowledges partial financial assistance by the Natural Sciences and Engineering Council of Canada.

a V -category X is its image under the Yoneda functor $y : X \rightarrow V^{X^{\text{op}}} = \hat{X}$, and its L -completion is the L -closure of that image in \hat{X} .

The main part of the paper is devoted to a substantial generalization of the first part which, however, without the reader's recalling of the more familiar V -category context, may be hard to motivate, especially in view of the considerable additional "technical" difficulties. The quantale V gets augmented by a *topological theory* $\mathcal{T} = (\mathbb{T}, V, \xi)$ which now entails also a Set -monad \mathbb{T} and a \mathbb{T} -algebra structure ξ on V , with suitable compatibility conditions (see [Hof07]). While a V -category X comes with a V -relation $a : X \multimap X$ (given by a function $a : X \times X \rightarrow V$), \mathcal{T} -categories come with a V -relation $a : TX \multimap X$ making X a lax \mathbb{T} -algebra. For \mathbb{T} the ultrafilter monad and $V = 2$, $\mathcal{T}\text{-Cat}$ provides Barr's [Bar70] relational description of the category of topological spaces (which, in turn, was based on Manes' [Man69] description of compact Hausdorff spaces); for the same monad but with V the Lawvere half-line, one obtains Lowen's approach spaces [Low89], as shown by Clementino and Hofmann [CH03].

The V -to- \mathcal{T} generalization must necessarily entail the provision of a Yoneda functor for a \mathcal{T} -category X . But what is X^{op} supposed to be in this highly asymmetric context? Fortunately, this problem was solved in [CH07]: the underlying set of X^{op} is TX , provided with a suitable \mathcal{T} -structure. This structure needs to be considered in addition to the free \mathbb{T} -algebra structure on TX , leading to the surprising fact that the \mathcal{T} -equivalent of the Yoneda functor of the familiar V -context has now two equally important facets. Once one has fully understood this "technical" part of the general theory, it is in fact rather straightforward to extend the V -categorical results on L -completion and L -separation to \mathcal{T} -categories, again with the help of the L -closure. We could therefore often keep the proofs in the \mathcal{T} -context quite short, especially when no new ideas beyond the initial "Yoneda investment" are needed.

Completeness of V -categories and the induced topology was also investigated by Flagg [Fla97, Fla92] (who called them V -continuity spaces). An alternative approach to the categories of interest in this paper was presented by Burroni [Bur71].

1. PRELIMINARIES

1.1. The quantale V . Throughout the paper we consider a commutative and unital quantale $V = (V, \otimes, k)$. Hence, V is a complete lattice with a commutative binary operation \otimes and neutral element k , such that $u \otimes (-)$ preserves suprema, for all $u \in V$. Consequently, V has an "internal hom" $u \multimap (-)$, given by

$$z \leq u \multimap v \iff z \otimes u \leq v,$$

for all $z, u, v \in V$. Sometimes we write $v \multimap u$ instead of $u \multimap v$. The quantale is *trivial* when $V = 1$; equivalently, when $k = \perp$ is the bottom element of V . Non-trivial examples of quantales are the two-element chain $2 = (\{0, 1\}, \wedge, 1)$, the extended positive half-line $P_+ = ([0, \infty]^{\text{op}}, +, 0)$, and $P_{\max} = ([0, \infty]^{\text{op}}, \max, 0)$; here $[0, \infty]^{\text{op}} = ([0, \infty], \geq)$, with the natural \geq . (We will use \vee, \wedge to denote suprema, infima in V , but use \sup, \inf, \max , etc. when we work in $[0, \infty]$ and refer to the natural order \leq .)

1.2. V -relations. The category $V\text{-Rel}$ has sets as objects, and a morphism $r : X \multimap Y$ is simply a function $r : X \times Y \rightarrow V$; its composite with $s : Y \multimap Z$ is given by

$$s \cdot r(x, z) = \bigvee_{y \in Y} r(x, y) \otimes s(y, z).$$

There is a functor

$$\text{Set} \longrightarrow V\text{-Rel}$$

which maps objects identically and interprets a map $f : X \rightarrow Y$ as a V -relation $f_\circ : X \leftrightarrow Y$:

$$f_\circ(x, y) = \begin{cases} k & \text{if } f(x) = y, \\ \perp & \text{otherwise;} \end{cases}$$

we normally write f instead of f_\circ . The functor is faithful precisely when $k > \perp$. The hom-sets of $V\text{-Rel}$ carry the pointwise order of V , so that $V\text{-Rel}$ becomes a 2-category. In fact, $V\text{-Rel}$ is Sup -enriched (with Sup the category of complete lattices and suprema-preserving maps), hence it is a quantaloid. Consequently, for every $r : X \leftrightarrow Y$, composition by r in $V\text{-Rel}$ from either side has a right adjoint, given by *extensions* and *liftings* respectively:

$$\begin{array}{ccc} (-) \cdot r \dashv (-) \bullet r & & r \cdot (-) \dashv r \bullet (-) \\ \frac{t \cdot r \leq s}{t \leq s \bullet r} & & \frac{r \cdot r \leq s}{t \leq r \bullet s} \\ \begin{array}{ccc} X & & \\ \downarrow r & \searrow s & \\ Y & \xrightarrow{t} & Z \end{array} & & \begin{array}{ccc} Y & & \\ \uparrow r & \swarrow s & \\ X & \xleftarrow{t} & Z \end{array} \\ s \bullet r(y, z) = \bigwedge_{x \in X} s(x, z) \circ r(x, y) & & r \bullet s(z, x) = \bigwedge_{y \in Y} r(x, y) \circ s(z, y) \end{array}$$

$V\text{-Rel}$ has a contravariant involution

$$(V\text{-Rel})^{\text{op}} \rightarrow V\text{-Rel}$$

which maps objects identically and assigns to $r : X \leftrightarrow Y$ its opposite relation $r^\circ : Y \leftrightarrow X$. When applied to a map $f = f_\circ$, one obtains $f \dashv f^\circ$ in the 2-category $V\text{-Rel}$.

1.3. V -categories. A V -category $X = (X, a)$ is a set X with a V -relation $a : X \leftrightarrow X$ satisfying $1_X \leq a$, $a \cdot a \leq a$; equivalently,

$$k \leq a(x, x), \quad a(x, y) \otimes a(y, z) \leq a(x, z)$$

for all $x, y, z \in X$. A V -functor $f : (X, a) \rightarrow (Y, b)$ must satisfy $f \cdot a \leq b \cdot f$; equivalently,

$$a(x, y) \leq b(f(x), f(y))$$

for all $x, y \in X$. The resulting category $V\text{-Cat}$ is the category Ord of (pre)ordered sets if $V = 2$, Lawvere's category Met of (pre)metric spaces if $V = \mathbb{P}_+$ (see [Law73]), and the category UMet of (pre)ultrametric spaces if $V = \mathbb{P}_{\max}$. For the trivial quantale one has $1\text{-Cat} = \text{Set}$. Furthermore, $V = (V, \circ)$ with its internal hom becomes a V -category.

$V\text{-Cat}$ is a symmetric monoidal closed category, with tensor product

$$(X, a) \otimes (Y, b) = (X \times Y, a \otimes b), \quad a \otimes b((x, y), (x', y')) = a(x, x') \otimes b(y, y'),$$

and internal hom

$$(X, a) \multimap (Y, b) = (V\text{-Cat}(X, Y), [a, b]), \quad [a, b](f, g) = \bigwedge_{x \in X} b(f(x), g(x)).$$

The \otimes -neutral object is $E = (E, k)$ (with a singleton set E), which generally must be distinguished from the terminal object $1 = (1, \top)$ in $V\text{-Cat}$. The internal hom describes the pointwise order if $V = 2$, and the usual sup-metric if $\mathbb{P}_+, \mathbb{P}_{\max}$.

1.4. **V-modules.** The category $\mathbf{V}\text{-Mod}$ has \mathbf{V} -categories as objects, and a morphism $\varphi : (X, a) \dashrightarrow (Y, b)$ is a \mathbf{V} -relation $\varphi : X \dashrightarrow Y$ with $\varphi \cdot a \leq \varphi$ and $b \cdot \varphi \leq \varphi$. Since always $\varphi = \varphi \cdot 1_X \leq \varphi \cdot a$ and $\varphi = 1_Y \cdot \varphi \leq b \cdot \varphi$, one actually has $\varphi \cdot a = \varphi$ and $b \cdot \varphi = \varphi$ for a \mathbf{V} -module $\varphi : X \dashrightarrow Y$. In particular, the \mathbf{V} -module $a : X \dashrightarrow X$ assumes the role of the identity morphism on X in $\mathbf{V}\text{-Mod}$, and we write $a = 1_X^*$, in order not to confuse it with 1_X in $\mathbf{V}\text{-Cat}$. This notation is extended to arbitrary maps $f : X \rightarrow Y$ by

$$f_* = b \cdot f, \quad f^* = f^\circ \cdot b,$$

and one easily verifies:

Lemma 1.1. *The following are equivalent for a map $f : X \rightarrow Y$ between \mathbf{V} -categories X and Y .*

- (i) $f : X \rightarrow Y$ is a \mathbf{V} -functor.
- (ii) f_* is a \mathbf{V} -module $f_* : X \dashrightarrow Y$.
- (iii) f^* is a \mathbf{V} -module $f^* : Y \dashrightarrow X$.

Hence there are functors which make the following diagram commute.

$$\begin{array}{ccccc} \mathbf{V}\text{-Cat} & \xrightarrow{(-)_*} & \mathbf{V}\text{-Mod} & \xleftarrow{(-)^*} & (\mathbf{V}\text{-Cat})^{\text{op}} \\ \uparrow & & \uparrow & & \uparrow \\ \mathbf{Set} & \xrightarrow{(-)_\circ} & \mathbf{V}\text{-Rel} & \xleftarrow{(-)^\circ} & \mathbf{Set}^{\text{op}} \end{array}$$

Here the vertical full embeddings are given by $X \mapsto (X, 1_X)$. Just like $\mathbf{V}\text{-Rel}$ also $\mathbf{V}\text{-Mod}$ is a quantaloid, with the same pointwise order structure. But not just suprema of \mathbf{V} -modules formed in $\mathbf{V}\text{-Rel}$ are again \mathbf{V} -modules, also extensions and liftings. For example, for $\varphi : (X, a) \dashrightarrow (Y, b)$, $\psi : (Z, c) \dashrightarrow (Y, b)$, the lifting $\varphi \dashrightarrow \psi$ formed in $\mathbf{V}\text{-Rel}$ is indeed a \mathbf{V} -module $\varphi \dashrightarrow \psi : (Z, c) \dashrightarrow (Y, b)$: from $\psi \cdot c \leq \psi$ and $\varphi \cdot (\varphi \dashrightarrow \psi) \leq \psi$ one obtains $\varphi \cdot (\varphi \dashrightarrow \psi) \cdot c \leq \psi$ and then $(\varphi \dashrightarrow \psi) \cdot c \leq \varphi \dashrightarrow \psi$; similar $a \cdot (\varphi \dashrightarrow \psi) \leq \varphi \dashrightarrow \psi$. Also the contravariant involution of $\mathbf{V}\text{-Rel}$ extends to $\mathbf{V}\text{-Mod}$ (e.g., if $\varphi : X \dashrightarrow Y$, then $\varphi : X^{\text{op}} \dashrightarrow Y^{\text{op}}$, where $X^{\text{op}} = (X, a^\circ)$ is the usual opposite \mathbf{V} -category.), and one has the commutative diagram.

$$\begin{array}{ccc} (\mathbf{V}\text{-Mod})^{\text{op}} & \xrightarrow{(-)^{\text{op}}} & \mathbf{V}\text{-Mod} \\ (-)^* \uparrow & & \uparrow (-)_* \\ \mathbf{V}\text{-Cat} & \xrightarrow{(-)^{\text{op}}} & \mathbf{V}\text{-Cat} \end{array}$$

As a quantaloid, $\mathbf{V}\text{-Mod}$ is in particular a 2-category, and for all $f : X \rightarrow Y$ in $\mathbf{V}\text{-Cat}$ one has

$$f_* \dashv f^*$$

in $\mathbf{V}\text{-Mod}$. $\mathbf{V}\text{-Cat}$ inherits its 2-categorical structure from $\mathbf{V}\text{-Mod}$ via

$$\begin{aligned} f \leq f' &: \iff f^* \leq (f')^* \iff \forall x \in X, y \in Y. b(y, f(x)) \leq b(y, f'(x)) \\ &\iff f'_* \leq f_* \iff \forall x \in X, y \in Y. b(f'(x), y) \leq b(f(x), y) \\ &\iff 1_X^* \leq (f')^* \cdot f_* \iff \forall x \in X. k \leq b(f(x), f'(x)). \end{aligned}$$

Hence, the previous diagram actually shows commuting 2-functors when we add dualization w.r.t. 2-cells (indicated by \circ) appropriately:

$$\begin{array}{ccc} (\mathbf{V}\text{-Mod})^{\text{op}} & \xrightarrow{(-)^{\text{op}}} & \mathbf{V}\text{-Mod} \\ (-)^* \uparrow & & \uparrow (-)_* \\ \mathbf{V}\text{-Cat} & \xrightarrow{(-)^{\text{op}}} & \mathbf{V}\text{-Cat}^{\text{co}} \end{array}$$

Of course, $\mathbf{V}\text{-Cat}$ being a 2-category, there is also a notion of adjointness in $\mathbf{V}\text{-Cat}$:

$$\begin{aligned}
f \dashv g \text{ in } \mathbf{V}\text{-Cat} &\iff f \cdot g \leq 1 \text{ and } 1 \leq g \cdot f \text{ in } \mathbf{V}\text{-Cat} \\
&\iff g^* \cdot f^* \leq 1^* \text{ and } 1^* \leq f^* \cdot g^* \text{ in } \mathbf{V}\text{-Mod} \\
&\iff g^* \dashv f^* \text{ in } \mathbf{V}\text{-Mod} \\
&\iff f_* = g^* \quad (\text{since } f_* \dashv f^* \text{ in } \mathbf{V}\text{-Mod}) \\
&\iff g_* = f_* \quad (\text{since } g_* \dashv g^* \text{ in } \mathbf{V}\text{-Mod}) \\
&\iff \forall x \in X, y \in Y. a(x, g(y)) = b(f(x), y).
\end{aligned}$$

1.5. **Yoneda.** \mathbf{V} -modules give rise to \mathbf{V} -functors, as follows

Proposition 1.2. *The following are equivalent for \mathbf{V} -relations $\varphi : X \dashv\vdash Y$ between \mathbf{V} -categories:*

- (i) $\varphi : X \dashv\vdash Y$ is a \mathbf{V} -module.
- (ii) $\varphi : X^{\text{op}} \otimes Y \longrightarrow \mathbf{V}$ is a \mathbf{V} -functor.

With $\varphi = a = 1_X^* : X \dashv\vdash X$ we obtain in particular the \mathbf{V} -functor $a : X^{\text{op}} \otimes X \longrightarrow \mathbf{V}$ whose transpose $\lceil a \rceil$ is the Yoneda- \mathbf{V} -functor

$$y : X \longrightarrow \hat{X} := (X^{\text{op}} \multimap \mathbf{V}), \quad x \longmapsto a(-, x).$$

The structure \hat{a} of \hat{X} is given by

$$\hat{a}(f, f') = \bigwedge_{x \in X} f(x) \multimap f'(x).$$

Lemma 1.3. *For all $x \in X$ and $f \in \hat{X}$, $\hat{a}(y(x), f) = f(x)$.*

One calls a \mathbf{V} -functor $f : (X, a) \longrightarrow (Y, b)$ *fully faithful* if $a(x, x') = b(f(x), f(x'))$ for all $x, x' \in X$; equivalently, if $1_X^* = f^* \cdot f_*$ (since $f^* \cdot f_* = f^\circ \cdot b \cdot f$), or just $1_X^* \geq f^* \cdot f_*$ (since the other inequality comes for free).

Corollary 1.4. *$y : X \longrightarrow \hat{X}$ is fully faithful.*

1.6. **L-separation.** For \mathbf{V} -functors $f, g : Z \longrightarrow X$ we write $f \cong g$ if $f \leq g$ and $g \leq f$; equivalently, if $f^* = g^*$, or $f_* = g_*$. We call X *L-separated* if $f \cong g$ implies $f = g$, for all $f, g : Z \longrightarrow X$. The full subcategory of $\mathbf{V}\text{-Cat}$ consisting of all L-separated \mathbf{V} -categories is denoted by $\mathbf{V}\text{-Cat}_{\text{sep}}$. Obviously, it suffices to consider $Z = E$ (the \otimes -neutral object) here: writing $x : E \longrightarrow X$ in $\mathbf{V}\text{-Cat}$ instead of $x \in X$, we just note that $f_* = g_*$ implies

$$(f \cdot x)_* = f_* \cdot x_* = g_* \cdot x_* = (g \cdot x)_*.$$

This proves the equivalence of (i),(ii) of the following proposition.

Proposition 1.5. *The following statements are equivalent for a \mathbf{V} -category $X = (X, a)$.*

- (i) X is L-separated.
- (ii) $x \cong y$ implies $x = y$, for all $x, y \in X$.
- (iii) For all $x, y \in X$, if $a(x, y) \geq k$ and $a(y, x) \geq k$, then $x = y$.
- (iv) The Yoneda functor $y : X \longrightarrow \hat{X}$ is injective.

Proof. For (ii) \iff (iii) \iff (iv) one observes

$$\begin{aligned} y(x) = y(y) &\iff x^\circ \cdot a = y^\circ \cdot a \\ &\iff x^* = y^* \\ &\iff x \leq y \text{ and } y \leq x \\ &\iff k \leq a(x, y) \text{ and } k \leq a(y, x). \end{aligned} \quad \square$$

Corollary 1.6. *The \mathbf{V} -category \mathbf{V} is L -separated. For all \mathbf{V} -categories X, Y , if Y is L -separated, $X \multimap Y$ is also L -separated. In particular, \hat{X} is L -separated, for every X .*

Proof. $k \leq u \multimap v$ and $k \leq v \multimap u$ means $u \leq v$ and $v \leq u$ in \mathbf{V} , hence $u = v$. For $Y = (Y, b)$ and $X = (X, a)$, $k \leq [a, b](f, g)$ in $X \multimap Y$ means $k \leq b(f(x), g(x))$ for all $x \in X$, which makes the second statement obvious. \square

1.7. L-completeness. Following Lawvere [Law73] we call a \mathbf{V} -category X *L-complete* if every adjunction $\varphi \dashv \psi : X \multimap Z$ in $\mathbf{V}\text{-Mod}$ is of the form $f_* \dashv f^*$, for a \mathbf{V} -functor $f : Z \rightarrow X$. Clearly, if X is L -separated, such a presentation is unique. As in 1.6, it suffices to consider $Z = E$ here; but we need the Axiom of Choice here.

Proposition 1.7. *The following statements are equivalent for a \mathbf{V} -category X .*

- (i) X is L -complete.
- (ii) Each left adjoint \mathbf{V} -module $\varphi : E \multimap X$ is of the form $\varphi = x_*$ for some x in X .
- (iii) Each right adjoint \mathbf{V} -module $\psi : X \multimap E$ is of the form $\psi = x^*$ for some x in X .

Elements in \hat{X} are \mathbf{V} -functors $X^{\text{op}} \cong X^{\text{op}} \otimes E \rightarrow \mathbf{V}$ which, by Proposition 1.2, may be considered as a \mathbf{V} -module $\psi : X \multimap E$. Suppose this \mathbf{V} -module has a left adjoint $\varphi : E \multimap X$. From $\varphi \cdot \psi \leq 1_X^*$ one obtains $\varphi \leq 1_X^* \bullet \psi$ (see 1.2), and from $(1_X^* \bullet \psi) \cdot \psi \leq 1_X^*$ and $\psi \cdot \varphi \geq 1_E^*$ one has $1_X^* \bullet \psi \leq \varphi$. Hence, if ψ is right adjoint, its left adjoint must necessarily be $1_X^* \bullet \psi$; moreover $(1_X^* \bullet \psi) \cdot \psi \leq 1_X^*$ always holds. Therefore:

Proposition 1.8. *A \mathbf{V} -module $\psi : X \multimap E$ (with $X = (X, a)$) is right adjoint if, and only if, $1_E^* \leq \psi \cdot (1_X^* \bullet \psi)$, that is, if*

$$(*) \quad k \leq \bigvee_{y \in Y} \psi(y) \otimes \left(\bigwedge_{x \in X} a(x, y) \multimap \psi(x) \right).$$

Note that $\bigwedge_{x \in X} a(x, y) \multimap \psi(x) = \hat{a}(\psi, y(y))$. We call a \mathbf{V} -functor $\psi : X^{\text{op}} \rightarrow \mathbf{V}$ *tight* if, as a \mathbf{V} -module $X \multimap E$, it is right adjoint, that is, if it satisfies (*). We consider

$$\tilde{X} = \{\psi \in \hat{X} \mid \psi \text{ tight}\}$$

as a full \mathbf{V} -subcategory of \hat{X} . Our goal is to exhibit \tilde{X} as an “ L -completion” of X .

Examples 1.9. (1) $\mathbf{V} = \mathbf{2}$. A \mathbf{V} -functor $X^{\text{op}} \rightarrow \mathbf{2}$ is the characteristic function of a down-closed set A in the (pre)ordered set X . Condition (*) the reads as

$$\exists y \in A \forall x \in A. x \leq y,$$

so that $A = \downarrow y$. In other words, \tilde{X} is simply the image of the Yoneda functor $y : X \rightarrow \hat{X}$, $y \mapsto \downarrow y$.

(2) $\mathbf{V} = \mathbf{P}_+$. A tight \mathbf{V} -functor $X^{\text{op}} \rightarrow \mathbf{V}$ is given by a function $\psi : X \rightarrow [0, \infty]$ with

$$\begin{aligned} \psi(y) \leq \psi(x) &\implies \psi(x) - \psi(y) \leq a(x, y) \quad (x, y \in X), \\ \inf_{y \in Y} (\psi(y) + \sup_{\substack{x \in X, \\ \psi(x) \leq a(x, y)}} (a(x, y) - \psi(x))) &= 0, \end{aligned}$$

here a is the generalized metric on X . If a is symmetric (so that $a = a^\circ$), these conditions are more conveniently describes as

$$|\psi(x) - \psi(y)| \leq a(x, y) \leq \psi(x) + \psi(y) \quad (x, y \in X),$$

$$\inf_{x \in X} \psi(x) = 0.$$

These are precisely the *supertight* maps on X considered in [LS00].

(3) $\mathbf{V} = \mathbf{P}_{\max}$. Here the two conditions of (2) change to

$$\psi(y) < \psi(x) \Rightarrow \psi(x) \leq a(x, y) \quad (x, y \in X),$$

$$\inf_{y \in Y} (\max(\psi(y), \sup_{\substack{x \in X, \\ \psi(x) < a(x, y)}} (a(x, y)))) = 0.$$

1.8. L-injectivity. A \mathbf{V} -functor $f : (X, a) \rightarrow (Y, b)$ is called *L-dense* if $f_* \cdot f^* = 1_Y^*$; that is, if $b = b \cdot f \cdot f^\circ \cdot b$, or

$$b(y, y') = \bigvee_{x \in X} b(y, f(x)) \otimes b(f(x), y')$$

for all $y, y' \in Y$. L-dense \mathbf{V} -functors have good composition-cancellation properties.

Lemma 1.10. *Let $f : X \rightarrow Y$ and $g : Y \rightarrow Z$ be \mathbf{V} -functors. Then the following assertions hold.*

- (1) f, g L-dense $\Rightarrow g \cdot f$ L-dense.
- (2) $g \cdot f$ L-dense $\Rightarrow g$ L-dense.
- (3) $g \cdot f$ L-dense, g fully faithful $\Rightarrow f$ L-dense.
- (4) $g \cdot f$ fully faithful, f L-dense $\Rightarrow g$ fully faithful.

A fully faithful L-dense \mathbf{V} -functor is an *L-equivalence*. Hence, f is an L-equivalence if, and only if, f_* (or f^*) is an isomorphism in $\mathbf{V}\text{-Mod}$. A \mathbf{V} -category Z is *pseudo-injective* if, for every fully faithful \mathbf{V} -functor $f : X \rightarrow Y$ and for all \mathbf{V} -functors $h : X \rightarrow Z$ there is a \mathbf{V} -functor $g : Y \rightarrow Z$ with $g \cdot f \cong h$; if strict equality is obtainable, we call Z *injective*. Z is *L-injective* if this extension property is required only for L-equivalences f . Hence, injectivity implies pseudo-injectivity, and every pseudo-injective \mathbf{V} -category is also L-injective.

Lemma 1.11. *The \mathbf{V} -category \mathbf{V} is injective, hence in particular L-injective.*

Proof. Let $f : X \rightarrow Y$ be fully faithful and $\varphi : X \rightarrow \mathbf{V}$ be any \mathbf{V} -functor. Then the \mathbf{V} -module $\varphi : E \rightarrow X$ factors as $\varphi = f^* \cdot \psi$, with $\psi = f_* \cdot \varphi$. But the \mathbf{V} -module $f^* \cdot \psi$ corresponds to the \mathbf{V} -functor $\psi \cdot f$, hence $\psi \cdot f = \varphi$.

$$\begin{array}{ccc} & \mathbf{V} & \\ \varphi \nearrow & \uparrow \psi & \\ X & \xrightarrow{f} & Y \end{array} \quad \begin{array}{ccc} & Y & \xrightarrow{f^*} & X \\ \psi = f_* \cdot \varphi \uparrow & \uparrow & \nearrow \varphi & \\ E & & & \end{array}$$

□

Note that the \mathbf{V} -functor ψ has been constructed effectively, with

$$\psi(y) = \bigvee_{x \in X} \varphi(x) \otimes b(f(x), y).$$

In case $\mathbf{V} = 2$, this means

$$\psi(y) = \top \iff \exists x \in X. (\varphi(x) = \top \text{ and } f(x) = y),$$

and for $V = \mathbf{P}_+$ we have

$$\psi(y) = \inf_{x \in X} (\varphi(x) + b(f(x), y)).$$

Proposition 1.12. *For all V -categories X, Y , if Y is pseudo-injective or L -injective, $X \multimap Y$ has the respective property. In particular, \hat{X} is injective.*

Proof. Let $f : A \rightarrow B$ be a fully faithful, and consider any V -functor $\varphi : A \rightarrow (X \multimap Y)$, with Y pseudo-injective. Since $f \otimes 1_X$ is fully faithful, the mate $\varphi_{\lrcorner} : A \otimes X \rightarrow Y$ factors (up to \cong) as $\varphi_{\lrcorner} \cong \psi_{\lrcorner} \cdot (f \otimes 1_X)$, with $\psi_{\lrcorner} : B \otimes X \rightarrow Y$ corresponding to a V -functor $\psi : B \rightarrow (X \multimap Y)$. Since $\psi_{\lrcorner} \cdot (f \otimes 1_X)$ corresponds to $\psi \cdot f$, $\varphi \cong \psi \cdot f$ follows. The proof works *mutatis mutandis* for L -injectivity. \square

Our goal is to show that L -injectivity and L -completeness are equivalent properties.

2. L-CLOSURE

2.1. L-dense V -functors. We first show that L -dense V -functors are characterized as “epimorphisms up to \cong ”.

Proposition 2.1. *A V -functor $m : M \rightarrow X$ is L -dense if, and only if, for all V -functors $f, g : X \rightarrow Y$ with $f \cdot m = g \cdot m$ one has $f \cong g$.*

Proof. The necessity of the condition is clear since from $f_* \cdot m_* = g_* \cdot m_*$ one obtains $f_* = g_*$ when $m_* \cdot m^* = 1_X^*$. To show the converse implication, by Lemma 1.10 we may assume that m is a full embedding $M \hookrightarrow X$ and consider its cokernel pair

$$(X, a) \begin{array}{c} \xrightarrow{f} \\ \xrightarrow{g} \end{array} (Y, b),$$

given by the disjoint union

$$Y = \{f(x) = g(x) \mid x \in M\} \cup \{f(x) \mid x \in X \setminus M\} \cup \{g(x) \mid x \in X \setminus M\},$$

where both f and g are full embeddings, and

$$b(f(x), g(y)) = \bigvee_{z \in Z} a(x, z) \otimes a(z, y)$$

for all $y, x \in X \setminus M$. Since $f_* = g_*$ by hypothesis, we obtain

$$a(x, y) = b(g(x), g(y)) = b(f(x), g(y)) = m_* \cdot m^*(x, y)$$

for all $x, y \in X \setminus M$. But this identity holds trivially when $x \in M$ or $y \in M$. Hence $m_* \cdot m^* = 1_X^*$. \square

Since $f \cong g$ precisely when $f \cdot x \cong g \cdot x$ for all $x \in X$ (considered as $x : E \rightarrow X$), it is now easy to identify the largest subset of X which contains M as an L -dense subset.

2.2. L-closure. For a V -category X and $M \subseteq X$, we define the L -closure of M in X by

$$\overline{M} = \{x \in X \mid \forall f, g : X \rightarrow Y. (f|_M = g|_M \Rightarrow f \cdot x \cong g \cdot x)\}$$

and prove

Proposition 2.2. *Let $X = (X, a)$ be a V -category, $M \subseteq X$ and $x \in X$. Then the following assertions are equivalent.*

- (i) $x \in \overline{M}$.
- (ii) $a(x, x) \leq \bigvee_{y \in M} a(x, y) \otimes a(y, x)$
- (iii) $k \leq \bigvee_{y \in M} a(x, y) \otimes a(y, x)$

- (iv) $1_E^* \leq x^* \cdot m_* \cdot m^* \cdot x_*$,
- (v) $m^* \cdot x_* \dashv x^* \cdot m_*$, where m denotes the full embedding $m : M \hookrightarrow X$.
- (vi) $x_* : E \dashv X$ factors through $m_* : M \dashv X$ by a map $\varphi : E \dashv M$ in $\mathbf{V}\text{-Mod}$.

Proof. (i) \Rightarrow (ii) follows from $M \hookrightarrow \overline{M}$ is dense. (ii) \Rightarrow (iii) is clear since $k \leq a(x, x)$. To see (iii) \Rightarrow (iv), just observe that

$$x^* \cdot m_* \cdot m^* \cdot x_*(\star, \star) = \bigvee_{y \in M} a(x, y) \otimes a(y, x).$$

Since $m^* \cdot x_* \cdot x^* \cdot m_* \leq m^* \cdot m_* = 1_M^*$, (iv) \Rightarrow (v). Assuming (v), we have $m_* \cdot m^* \cdot x_* \dashv x^* \cdot m_* \cdot x^*$ as well as $m_* \cdot m^* \cdot x_* \leq x_*$ and $x^* \cdot m_* \cdot x^* \leq x^*$, which implies $m_* \cdot m^* \cdot x_* = x_*$ and we have shown (v) \Rightarrow (vi). Finally, assume (vi) and let $f, g : X \rightarrow Y$ with $f|_M = g|_M$. Then

$$f_* \cdot x_* = f_* \cdot m_* \cdot \varphi = g_* \cdot m_* \cdot \varphi = g_* \cdot x_*,$$

which proves (i). □

V-functors respect the L-closure, as we show next.

Proposition 2.3. For a \mathcal{J} -functor $f : X \rightarrow Y$ and $M, M' \subseteq X, N \subseteq Y$, we have

- (1) $M \subseteq \overline{M}$ and $M \subseteq M'$ implies $\overline{M} \subseteq \overline{M'}$.
- (2) $\overline{\emptyset} = \emptyset$ and $\overline{\overline{M}} = \overline{M}$.
- (3) $f(\overline{M}) \subseteq \overline{f(M)}$ and $f^{-1}(\overline{N}) \supseteq \overline{f^{-1}(N)}$.
- (4) If k is \vee -irreducible (so that $k \leq u \vee v$ implies $k \leq u$ or $k \leq v$), then $\overline{M \cup M'} = \overline{M} \cup \overline{M'}$.

Proof. (1), (2) are obvious. For (3), applying Lemma 1.10 to

$$\begin{array}{ccc} M & \longrightarrow & f(M) \\ \downarrow & & \downarrow \\ \overline{M} & \longrightarrow & f(\overline{M}) \end{array}$$

one sees that $f(M) \rightarrow f(\overline{M})$ is L-dense, hence $f(\overline{M}) \subseteq \overline{f(M)}$. With $M = f^{-1}(N)$, this implies $\overline{f^{-1}(N)} \subseteq f^{-1}(\overline{N})$. To see (4), we just need to show that $x \in \overline{M \cup M'}$ implies $x \in \overline{M}$ or $x \in \overline{M'}$. But this follows from

$$k \leq \bigvee_{y \in M \cup M'} a(x, y) \otimes a(y, x) = \left(\bigvee_{y \in M} a(x, y) \otimes a(y, x) \right) \vee \left(\bigvee_{y \in M'} a(x, y) \otimes a(y, x) \right),$$

assuming that k is \vee -irreducible. □

Corollary 2.4. If k is \vee -irreducible in \mathbf{V} , then the L-closure operator defines a topology on X such that every V-functor becomes continuous. Hence, L-closure defines a functor $L : \mathbf{V}\text{-Cat} \rightarrow \mathbf{Top}$.

Examples 2.5. (1) For $X = (X, \leq)$ in $2\text{-Cat} = \mathbf{Ord}$ and $M \subseteq X$, one has $x \in \overline{M}$ precisely when $x \leq z \leq x$ for some $z \in M$. Also, $M \subseteq X$ is open in LX if every $x \in M$ satisfies

$$\forall z \in X. (x \leq z \leq x \Rightarrow z \in M).$$

(2) In \mathbf{Met} , $\overline{M} = \{x \in X = (X, a) \mid \inf_{z \in M} (a(x, z) + a(z, x)) = 0\}$, and in \mathbf{UMet}

$$\overline{M} = \{x \in X = (X, a) \mid \inf_{z \in M} (\max(a(x, z), a(z, x))) = 0\}$$

which for symmetric (ultra)metric spaces describes the ordinary topological closure.

2.3. L-separatedness via the L-closure.

Proposition 2.6. *Let $X = (X, a)$ be a \mathbf{V} -category and $\Delta \subseteq X \times X$ the diagonal. Then*

$$\bar{\Delta} = \{(x, y) \in X \times Y \mid x \cong y\}$$

Proof. Let first $(x, y) \in \bar{\Delta}$. With $\pi_1, \pi_2 : X \times X \rightarrow X$ denoting the projection maps, we have $\pi_1|_{\Delta} = \pi_2|_{\Delta}$ and therefore $x = \pi_1(x, y) \cong \pi_2(x, y) = y$. Assume now $x \cong y$. Note that the canonical functor $\mathbf{V}\text{-Cat} \rightarrow \mathbf{Ord}$ preserves products, hence

$$(x_1, y_1) \cong (x_2, y_2) \iff x_1 \cong x_2 \text{ and } y_1 \cong y_2,$$

for all $(x_1, y_1), (x_2, y_2) \in X \times X$. Therefore we have $(x, y) \cong (x, x)$. Let now $f, g : X \times X \rightarrow Y$ be \mathbf{V} -functors with $f|_{\Delta} = g|_{\Delta}$. Then $f(x, y) \cong f(x, x) = g(x, x) \cong g(x, y)$. \square

Corollary 2.7. *A \mathbf{V} -category X is L-separated if and only if the diagonal Δ is closed in $X \times X$.*

Theorem 2.8. *$\mathbf{V}\text{-Cat}_{\text{sep}}$ is an epi-reflective subcategory of $\mathbf{V}\text{-Cat}$, where the reflection map is given by $y_X : X \rightarrow y_X(X)$, for each \mathbf{V} -category X . Hence, limits of L-separated \mathbf{V} -categories are formed in $\mathbf{V}\text{-Cat}$, while colimits are obtained by reflecting the colimit formed in $\mathbf{V}\text{-Cat}$. The epimorphisms in $\mathbf{V}\text{-Cat}$ are precisely the L-dense \mathbf{V} -functors.*

2.4. L-completeness via the L-closure.

Lemma 2.9. *Let $X = (X, a)$ be a \mathbf{V} -category and $M \subseteq X$. Then the following assertions hold.*

- (1) *Assume that X is L-complete and M be L-closed. Then M is L-complete.*
- (2) *Assume that X is L-separated and M is L-complete. Then M is L-closed.*

Proof. (1) follows immediately from Proposition 2.2. To see (2), let $x \in X$ such that $m^* \cdot x_* + x^* \cdot m_*$. Since M is L-complete, there is some $y \in M$ such that $y_* = m^* \cdot x_*$ and $y^* = x^* \cdot m_*$. Hence $m(y)_* = m_* \cdot y_* \leq x_*$ and $m(y)^* = y^* \cdot m^* \leq x^*$ and therefore, $m(y)_* = x_*$. L-separatedness of X gives now $m(y) = x$, i.e. $x \in M$. \square

Theorem 2.10. *Let $X = (X, b)$ be a \mathbf{V} -category. The following assertions are equivalent.*

- (i) *X is L-complete.*
- (ii) *X is L-injective.*
- (iii) *$y : X \rightarrow \tilde{X}$ has a pseudo left-inverse \mathbf{V} -functor $R : \tilde{X} \rightarrow X$, i.e. $R \cdot y \cong 1_X$.*

Proof. As for Theorem 4.11. \square

Proposition 2.11. *For a \mathbf{V} -category X , as a set \tilde{X} (see 1.7) coincides with the L-closure of $y(X)$ in \hat{X} . Hence, $y : X \rightarrow \tilde{X}$ is fully faithful and L-dense, and \tilde{X} is L-complete.*

Proof. By Proposition 2.2, a \mathbf{V} -functor $\psi : X^{\text{op}} \rightarrow \mathbf{V}$ lies in the L-closure of $y(X)$ in \hat{X} if, and only if,

$$k \leq \bigvee_{y \in X} \hat{a}(\psi, y(y)) \otimes \hat{a}(y(y), \psi).$$

Since $\hat{a}(y(y), \psi) = \psi(y)$ by Lemma 1.3, this means precisely that ψ must be tight. \square

Theorem 2.12. *The full subcategory $\mathbf{V}\text{-Cat}_{\text{cpl}}$ of $\mathbf{V}\text{-Cat}_{\text{sep}}$ of L-complete \mathbf{V} -categories is an epi-reflective subcategory of $\mathbf{V}\text{-Cat}_{\text{sep}}$. The reflection map of a L-separated \mathbf{V} -category X is given by any dense embedding of X into a L-complete and L-separated \mathbf{V} -category, for instance by $y : X \rightarrow \tilde{X}$.*

3. THE \mathcal{T} -SETTING

3.1. **The theory \mathcal{T} .** From now on we assume that the quantale \mathbf{V} is part of a *strict topological theory* $\mathcal{T} = (\mathbb{T}, \mathbf{V}, \xi)$ as introduced in [Hof07]. Here $\mathbb{T} = (T, e, m)$ is a **Set**-monad where T and m satisfy (BC) (that is, T sends pullbacks to weak pullbacks and each naturality square of m is a weak pullback) and $\xi : T\mathbf{V} \rightarrow \mathbf{V}$ is a map such that

$$1_{\mathbf{V}} = \xi \cdot e_{\mathbf{V}}, \quad \xi \cdot T\xi = \xi \cdot m_{\mathbf{V}},$$

the diagrams

$$\begin{array}{ccc} T(\mathbf{V} \times \mathbf{V}) & \xrightarrow{T(\otimes)} & T\mathbf{V} \\ \langle \xi \cdot T\pi_1, \xi \cdot T\pi_2 \rangle \downarrow & & \downarrow \xi \\ \mathbf{V} \times \mathbf{V} & \xrightarrow{\otimes} & \mathbf{V}, \end{array} \quad \begin{array}{ccc} T1 & \xrightarrow{Tk} & T\mathbf{V} \\ ! \downarrow & & \downarrow \xi \\ 1 & \xrightarrow{k} & \mathbf{V}, \end{array}$$

commute and

$(\xi_X)_X : P_{\mathbf{V}} \rightarrow P_{\mathbf{V}}T$ is a natural transformation, where $P_{\mathbf{V}}$ is the \mathbf{V} -powerset functor considered as a functor from **Set** to **Ord** and the X -component $\xi_X : P_{\mathbf{V}}(X) \rightarrow P_{\mathbf{V}}T(X)$ is given by $\varphi \mapsto \xi \cdot T\varphi$.

Here $P_{\mathbf{V}}(X) = \mathbf{V}^X$, and for a function $f : X \rightarrow Y$ we have a canonical map $f^{-1} : \mathbf{V}^Y \rightarrow \mathbf{V}^X$, $\varphi \mapsto \varphi \cdot f$. Now $P_{\mathbf{V}}(f)$ is defined as the left adjoint to f^{-1} , explicitly, for $\varphi \in \mathbf{V}^X$ we have $P_{\mathbf{V}}(\varphi)(y) = \bigvee_{x \in f^{-1}(y)} \varphi(x)$. Furthermore, we assume $T1 = 1$.

Examples 3.1. (1) For each quantale \mathbf{V} , $(\mathbb{1}, \mathbf{V}, 1_{\mathbf{V}})$ is a strict topological theory, where $\mathbb{1} = (\text{Id}, 1, 1)$ denotes the identity monad.

(2) $\mathcal{U}_2 = (\mathbb{U}, \mathbf{2}, \xi_2)$ is a strict topological theory, where $\mathbb{U} = (U, e, m)$ denotes the ultrafilter monad and ξ_2 is essentially the identity map.

(3) $\mathcal{U}_{\mathbf{P}_+} = (\mathbb{U}, \mathbf{P}_+, \xi_{\mathbf{P}_+})$ is a strict topological theory, where

$$\xi_{\mathbf{P}_+} : UP_+ \rightarrow \mathbf{P}_+, \quad x \mapsto \inf\{v \in \mathbf{P}_+ \mid [0, v] \in x\}.$$

As shown in [Hof07, Lemma 3.2], the right adjoint \dashv of the tensor product \otimes in \mathbf{V} is automatically compatible with the map $\xi : T\mathbf{V} \rightarrow \mathbf{V}$ in the sense that

$$\xi \cdot T(\dashv) \leq \dashv \cdot \langle \xi \cdot T\pi_1, \xi \cdot T\pi_2 \rangle.$$

$$\begin{array}{ccc} T(\mathbf{V} \times \mathbf{V}) & \xrightarrow{T(\dashv)} & T\mathbf{V} \\ \langle \xi \cdot T\pi_1, \xi \cdot T\pi_2 \rangle \downarrow & \geq & \downarrow \xi \\ \mathbf{V} \times \mathbf{V} & \xrightarrow{\dashv} & \mathbf{V} \end{array}$$

Furthermore, our condition $T1 = 1$ implies $m_X^\circ \cdot e_X = e_{TX} \cdot e_X$ for each set X . In fact, $m_X^\circ \cdot e_X \geq e_{TX} \cdot e_X$ is true for each monad since $m_X^\circ \geq e_{TX}$. Let now $\mathfrak{X} \in TTX$ and $x \in X$ such that $m_X(\mathfrak{X}) = e_X(x)$. We consider the commutative diagram

$$\begin{array}{ccc} TT1 & \xrightarrow{TTx} & TTX \\ m_1 \downarrow & & \downarrow m_X \\ T1 & \xrightarrow{Tx} & TX, \end{array}$$

where $x : 1 \mapsto X$. Since m satisfies (BC), there is some $\mathfrak{Y} \in TT1 = 1$ with $TTx(\mathfrak{Y}) = \mathfrak{X}$, that is, $\mathfrak{X} = e_{TX} \cdot e_X(x)$.

The functor $T : \mathbf{Set} \rightarrow \mathbf{Set}$ can be extended to a 2-functor $T_\xi : \mathbf{V}\text{-Rel} \rightarrow \mathbf{V}\text{-Rel}$ as follows. Given a \mathbf{V} -relation $r : X \times Y \rightarrow \mathbf{V}$, we define $T_\xi r : TX \times TY \rightarrow \mathbf{V}$ as the left Kan-extension

$$\begin{array}{ccc} T(X \times Y) & \xrightarrow{\text{can}} & TX \times TY \\ & \searrow \xi_{X \times Y}(r) = \xi \cdot Tr & \swarrow T_\xi r \\ & & \mathbf{V} \end{array}$$

in \mathbf{Ord} (where $TX, TY, T(X \times Y)$ are discrete), i.e. the smallest (order-preserving) map $s : TX \times TY \rightarrow \mathbf{V}$ such that $\xi \cdot Tr \leq s \cdot \text{can}$. Elementwise, we have

$$T_\xi r(\mathfrak{x}, \mathfrak{y}) = \bigvee \left\{ \xi \cdot Tr(w) \mid w \in T(X \times Y), T\pi_1(w) = \mathfrak{x}, T\pi_2(w) = \mathfrak{y} \right\}$$

for each $\mathfrak{x} \in TX$ and $\mathfrak{y} \in TY$. We have the following properties.

Proposition 3.2 ([Hof07]). *The following assertions hold.*

- (1) For each \mathbf{V} -matrix $r : X \leftrightarrow Y$, $T_\xi(r^\circ) = T_\xi(r)^\circ$ (and we write $T_\xi r^\circ$).
- (2) For each function $f : X \rightarrow Y$, $Tf = T_\xi f$ (and therefore also $Tf^\circ = T_\xi f^\circ$).
- (3) $e_Y \cdot r \leq T_\xi r \cdot e_X$ for all $r : X \leftrightarrow Y$ in $\mathbf{V}\text{-Rel}$.
- (4) $m_Y \cdot T_\xi^2 r = T_\xi r \cdot m_X$ for all $r : X \leftrightarrow Y$ in $\mathbf{V}\text{-Rel}$.

3.2. \mathcal{T} -relations. We define a \mathcal{T} -relation from X to Y to be a \mathbf{V} -relation of the form $a : TX \leftrightarrow Y$, and write $a : X \leftrightarrow Y$. Given also $b : Y \leftrightarrow Z$, the composite $b \circ a : X \leftrightarrow Z$ is given by the Kleisli convolution

$$b \circ a = b \cdot T_\xi a \cdot m_X^\circ.$$

Composition of \mathcal{T} -relations is associative, and for each \mathcal{T} -matrix $a : X \leftrightarrow Y$ we have $a \circ e_X^\circ = a$ and $e_Y^\circ \circ a \geq a$, hence $e_X^\circ : X \leftrightarrow X$ is a lax identity. We call a \mathcal{T} -relation $a : X \leftrightarrow Y$ *unitary* if $e_Y^\circ \circ a = a$, so that $e_X^\circ : X \leftrightarrow X$ is the identity on X in the category $\mathcal{T}\text{-URel}$ of sets and unitary \mathcal{T} -relations, with the Kleisli convolution as composition. The hom-sets of $\mathcal{T}\text{-URel}$ inherit the order-structure from $\mathbf{V}\text{-Rel}$, and composition of (unitary) \mathcal{T} -relations respects this order in both variables. Many notions and arguments can be transported from the \mathbf{V} -setting to the \mathcal{T} -setting by substituting relational composition by Kleisli convolution.

Given a \mathcal{T} -relation $c : X \leftrightarrow Z$, the composition by c from the right side has a right adjoint but composition by c from the left side in general not. Explicitely, given also $b : X \leftrightarrow Y$, we pass from

$$\begin{array}{ccc} X \xrightarrow{b} Y & \text{to} & TX \xrightarrow{b} Y \\ c \downarrow & & m_X^\circ \downarrow \\ Z & & TT X \\ & & T_\xi c \downarrow \\ & & TZ \end{array}$$

in $\mathcal{T}\text{-URel}$ in $\mathbf{V}\text{-Rel}$

and define the extension $b \circ\text{-} c : Z \leftrightarrow Y$ as $b \bullet (T_\xi c \cdot m_X^\circ) : TZ \leftrightarrow Y$.

3.3. \mathcal{T} -categories. A \mathcal{T} -category $X = (X, a)$ is a set X equipped with a \mathcal{T} -relation $a : X \leftrightarrow X$ satisfying $e_X^\circ \leq a$ and $a \circ a \leq a$, equivalently,

$$k \leq a(e_X(x), x), \quad T_\xi a(\mathfrak{X}, \mathfrak{x}) \otimes a(x, x) \leq a(x, x)$$

for all $\mathfrak{x} \in TTX$, $x \in TX$ and $x \in X$. A \mathcal{T} -functor $f : (X, a) \longrightarrow (Y, b)$ must satisfy $f \cdot a \leq b \cdot Tf$, which in pointwise notation reads as

$$a(\mathfrak{x}, x) \leq b(Tf(\mathfrak{x}), f(x))$$

for all $\mathfrak{x} \in TX$ and $x \in X$. The resulting category of \mathcal{T} -categories and \mathcal{T} -functors is denoted by $\mathcal{T}\text{-Cat}$ (see also [CH03, CT03, CHT04]). Note that the quantale \mathbf{V} becomes in a natural way a \mathcal{T} -category $\mathbf{V} = (\mathbf{V}, \text{hom}_\xi)$ where $\text{hom}_\xi : TV \times \mathbf{V} \longrightarrow \mathbf{V}$, $(v, v) \longmapsto (\xi(v) \multimap v)$.

Examples 3.3. (1) For each quantale \mathbf{V} , $\mathcal{J}_\mathbf{V}$ -categories are precisely \mathbf{V} -categories and $\mathcal{J}_\mathbf{V}$ -functors are \mathbf{V} -functors.

(2) The main result of [Bar70] states that $\mathcal{U}_2\text{-Cat}$ is isomorphic to the category Top of topological spaces. The \mathcal{U}_2 -category $\mathbf{V} = \mathbf{2}$ is the Sierpinski space with $\{0\}$ open and $\{1\}$ closed. In [CH03] it is shown that $\mathcal{U}_{\mathbf{P}_+}\text{-Cat}$ is isomorphic to the category App of approach spaces (see [Low97] for more details about App).

A \mathcal{T} -category $X = (X, a)$ can be also thought of as a lax Eilenberg–Moore algebra, since the two conditions above can be equivalently expressed as

$$1_X \leq a \cdot e_X, \quad a \cdot T_\xi a \leq a \cdot m_X.$$

As a consequence, each \mathbb{T} -algebra (X, a) can be considered as a \mathcal{T} -category by simply regarding the function $\alpha : TX \longrightarrow X$ as a \mathcal{T} -relation $\alpha : X \dashrightarrow X$. The free Eilenberg–Moore algebra (TX, m_X) – viewed as a \mathcal{T} -category – is denoted by $|X|$.

Each \mathcal{T} -category $X = (X, a)$ has an underlying \mathbf{V} -category $SX = (X, a \cdot e_X)$. Indeed, this defines a functor $S : \mathcal{T}\text{-Cat} \longrightarrow \mathbf{V}\text{-Cat}$ which has a left adjoint $A : \mathbf{V}\text{-Cat} \longrightarrow \mathcal{T}\text{-Cat}$ defined by $AX = (X, e_X^\circ \cdot T_\xi r)$, for each \mathbf{V} -category $X = (X, r)$. There is yet another interesting functor connecting \mathcal{T} -categories and \mathbf{V} -categories, namely $M : \mathcal{T}\text{-Cat} \longrightarrow \mathbf{V}\text{-Cat}$ which sends a \mathcal{T} -category (X, a) to the \mathbf{V} -category $(TX, T_\xi a \cdot m_X^\circ)$. The *dual \mathcal{T} -category* X^{op} (see [CH07]) of a \mathcal{T} -category $X = (X, a)$ is then defined as

$$X^{\text{op}} = A(M(X)^{\text{op}}).$$

Examples 3.4. For $\mathbb{T} = \mathbb{U}$ the ultrafilter monad, the topology on $|X|$ can be described via the Zariski-closure:

$$\mathfrak{x} \in \text{cl}(\mathcal{A}) \iff \mathfrak{x} \supseteq \bigcap \mathcal{A} \iff \bigcup \mathcal{A} \subseteq \mathfrak{x},$$

for $\mathfrak{x} \in UX$ and $\mathcal{A} \subseteq UX$. Furthermore, for $X \in \mathcal{U}_2\text{-Cat} \cong \text{Top}$, $M(X) = (UX, \leq)$ is the (pre)ordered set where

$$\mathfrak{x} \leq \mathfrak{y} \iff \forall A \in \mathfrak{x}. \bar{A} \in \mathfrak{y}$$

for $\mathfrak{x}, \mathfrak{y} \in UX$. Then X^{op} is the Alexandroff space induced by the dual order \geq . If $X \in \mathcal{U}_{\mathbf{P}_+}\text{-Cat} \cong \text{App}$ is an approach space with distance function $\text{dist} : PX \times X \longrightarrow \mathbf{P}_+$, then $M(X) = (UX, d)$ is the (generalized) metric space with

$$d(\mathfrak{x}, \mathfrak{y}) = \inf\{\varepsilon \in [0, \infty] \mid \forall A \in \mathfrak{x}. \bar{A}^{(\varepsilon)} \in \mathfrak{y}\},$$

where $\mathfrak{x}, \mathfrak{y} \in UX$ and $\bar{A}^{(\varepsilon)} = \{x \in X \mid \text{dist}(A, x) \leq \varepsilon\}$.

The tensor product of \mathbf{V} can be transported to $\mathcal{T}\text{-Cat}$ by putting $(X, a) \otimes (Y, b) = (X \times Y, c)$ with

$$c(w, (x, y)) = a(x, x) \otimes b(y, y),$$

where $w \in T(X \times Y)$, $x \in X$, $y \in Y$, $\mathfrak{x} = T\pi_1(w)$ and $\mathfrak{y} = T\pi_2(w)$. The \mathcal{T} -category $E = (E, k)$ is a \otimes -neutral object, where E is a singleton set and k the constant relation with value $k \in \mathbf{V}$. Unlike the \mathbf{V} -case, this does not result in general in a closed structure on $\mathcal{T}\text{-Cat}$. However, as shown in [Hof07], if a \mathcal{T} -category

$X = (X, a)$ satisfies $a \cdot T_\xi a = a \cdot m_X$, then $X \otimes_- : \mathcal{T}\text{-Cat} \rightarrow \mathcal{T}\text{-Cat}$ has a right adjoint ${}_X : \mathcal{T}\text{-Cat} \rightarrow \mathcal{T}\text{-Cat}$. Explicitly, for a \mathcal{T} -category $Y = (Y, b)$, the exponential $X \multimap Y$ is given by the set

$$\{f : X \rightarrow Y \mid f \text{ is a } \mathcal{T}\text{-functor}\}$$

equipped with the structure-relation $\llbracket a, b \rrbracket$ defined as

$$\llbracket a, b \rrbracket(\mathfrak{p}, h) = \bigvee \left\{ v \in \mathbf{V} \mid \forall q \in T\pi_2^{-1}(\mathfrak{p}), x \in X. a(T\pi_1(q), x) \otimes v \leq b(\text{Ev}(q), h(x)) \right\},$$

where $\mathfrak{p} \in T(Y^X)$, $h \in Y^X$, $\pi_1 : X \times (X \multimap Y) \rightarrow X$ and $\pi_2 : X \times (X \multimap Y) \rightarrow Y^X$. Using the adjunction $u \otimes_- \dashv u \multimap_-$ in \mathbf{V} , we see that

$$\llbracket a, b \rrbracket(\mathfrak{p}, h) = \bigwedge_{\substack{q \in T(X \times (X \multimap Y)), x \in X \\ q \mapsto \mathfrak{p}}} a(T\pi_1(q), x) \multimap b(\text{Ev}(q), h(x)).$$

Lemma 3.5. *Let $X = (X, a)$, $Y = (Y, b)$ be \mathcal{T} -categories with $a \cdot T_\xi a = a \cdot m_X$ and $h, h' \in (X \multimap Y)$. Then*

$$\llbracket a, b \rrbracket(e_{Y^X}(h), h') = \bigwedge_{x \in X} b(e_Y(h(x)), h'(x)).$$

3.4. \mathcal{T} -modules. Let $X = (X, a)$ and $Y = (Y, b)$ be \mathcal{T} -categories and $\varphi : X \multimap Y$ be a \mathcal{T} -relation. We call φ a \mathcal{T} -module, and write $\varphi : X \multimap Y$, if $\varphi \circ a \leq \varphi$ and $b \circ \varphi \leq \varphi$. Note that we have always $\varphi \circ a \geq \varphi$ and $b \circ \varphi \geq \varphi$, so that the \mathcal{T} -module condition above implies equality. It is easy to see that the extension as well as the lifting (if it exists) in $\mathcal{T}\text{-URel}$ of \mathcal{T} -modules is again a \mathcal{T} -module. Furthermore, we have $a : X \multimap X$ for each \mathcal{T} -category $X = (X, a)$; in fact, a is the identity \mathcal{T} -module on X for the Kleisli convolution. The category of \mathcal{T} -categories and \mathcal{T} -modules, with Kleisli convolution as composition is denoted by $\mathcal{T}\text{-Mod}$. In fact, $\mathcal{T}\text{-Mod}$ is an ordered category, with the structure on hom-sets inherited from $\mathcal{T}\text{-URel}$.

Let now $X = (X, a)$ and $Y = (Y, b)$ be \mathcal{T} -categories and $f : X \rightarrow Y$ be a Set -map. We define \mathcal{T} -relations $f_* : X \multimap Y$ and $f^* : Y \multimap X$ by putting $f_* = b \cdot T f$ and $f^* = f^\circ \cdot b$ respectively. Hence, for $\mathfrak{x} \in TX$, $\mathfrak{y} \in TY$, $x \in X$ and $y \in Y$, $f_*(\mathfrak{x}, \mathfrak{y}) = b(Tf(\mathfrak{x}), \mathfrak{y})$ and $f^*(\mathfrak{y}, \mathfrak{x}) = b(\mathfrak{y}, f(x))$. Given now \mathcal{T} -modules φ and ψ , we have

$$\varphi \circ f_* = \varphi \cdot T f \quad \text{and} \quad f^* \circ \psi = f^\circ \cdot \psi.$$

The latter equality follows from

$$f^* \circ \psi = f^\circ \cdot b \cdot T_\xi \psi \cdot m_Z^\circ = f^\circ \cdot \psi,$$

whereby the first equality follows from

$$\varphi \circ f_* = \varphi \circ (b \cdot T f) = \varphi \cdot T_\xi b \cdot T^2 f \cdot m_X^\circ = \varphi \cdot T_\xi b \cdot m_Y^\circ \cdot T f = \varphi \cdot T f.$$

In particular we have $b \circ f_* = f_*$ and $f^* \circ b = f^*$, as well as $f_* \circ f^* = b \cdot T f \cdot T f^\circ \cdot T_\xi b \cdot m_Y^\circ \leq b$. In the latter case we have even equality provided that f is surjective. As before, one easily verifies

Proposition 3.6. *The following assertions are equivalent.*

- (i) $f : X \rightarrow Y$ is a \mathcal{T} -functor.
- (ii) f_* is a \mathcal{T} -module $f_* : X \multimap Y$.
- (iii) f^* is a \mathcal{T} -module $f^* : Y \multimap X$.

As in the \mathbf{V} -case, we have functors

$$\mathcal{T}\text{-Cat} \xrightarrow{(-)_*} \mathcal{T}\text{-Mod} \xleftarrow{(-)^*} \mathcal{T}\text{-Cat}^{\text{op}}.$$

We can transport the order-structure on hom-sets from $\mathcal{T}\text{-Mod}$ to $\mathcal{T}\text{-Cat}$ via the functor $(-)^* : \mathcal{T}\text{-Cat}^{\text{op}} \rightarrow \mathcal{T}\text{-Mod}$, that is, we define $f \leq g$ if $f^* \leq g^*$, or equivalently, if $g_* \leq f_*$. With this definition we turn $\mathcal{T}\text{-Cat}$ into an *ordered category*. As usual, we call \mathcal{T} -functors $f, g : X \rightarrow Y$ *equivalent*, and write $f \cong g$, if $f \leq g$ and $g \leq f$. Hence, $f \cong g$ if and only if $f^* = g^*$, which in turn is equivalent to $f_* = g_*$.

Lemma 3.7. *Let $f, g : X \rightarrow Y$ be \mathcal{T} -functors between \mathcal{T} -categories $X = (X, a)$ and $Y = (Y, b)$. Then*

$$f \leq g \iff \forall x \in X. k \leq b(e_Y(f(x)), g(x)).$$

Proof. If $g_* \leq f_*$, then

$$k \leq g_*(e_X(x), g(x)) \leq f_*(e_X(x), g(x)) = b(e_Y(f(x)), g(x)).$$

On the other hand, if $k \leq b(e_Y(g(x)), f(x))$ for each $x \in X$, then

$$f^*(\eta, x) = b(\eta, f(x)) \leq T_{\xi} b(e_{TY}(\eta), e_Y(f(x))) \otimes b(e_Y(f(x)), g(x)) \leq b(\eta, g(x)) = g^*(\eta, x). \quad \square$$

In particular, for \mathcal{T} -functors $f, g : X \rightarrow V$, we have $f \leq g$ if and only if $f(x) \leq g(x)$ for all $x \in X$. Assume now that $X = (X, a)$, $Y = (Y, b)$ and $Z = (Z, c)$ are \mathcal{T} -categories where $a \cdot T_{\xi} a = a \cdot m_X$. By combining the previous lemma with Lemma 3.5, we obtain $f \leq g \iff \lceil f \rceil \leq \lceil g \rceil$ for all \mathcal{T} -functors $f, g : X \otimes Y \rightarrow Z$, where $\lceil f \rceil, \lceil g \rceil : Y \rightarrow Z^X$.

3.5. Yoneda. Also \mathcal{T} -modules give rise to \mathcal{T} -functors, but besides X^{op} we must take also the \mathcal{T} -category $|X|$ (see 3.3) into consideration.

Theorem 3.8 ([CH07]). *For \mathcal{T} -categories (X, a) and (Y, b) , and a \mathcal{T} -relation $\psi : X \dashv\dashv Y$, the following assertions are equivalent.*

- (i) $\psi : (X, a) \dashv\dashv (Y, b)$ is a \mathcal{T} -module.
- (ii) Both $\psi : |X| \otimes Y \rightarrow V$ and $\psi : X^{\text{op}} \otimes Y \rightarrow V$ are \mathcal{T} -functors.

Since we have $a : X \dashv\dashv X$ for each \mathcal{T} -category $X = (X, a)$, the theorem above provides us with two \mathcal{T} -functors

$$a : |X| \otimes X \rightarrow V \quad \text{and} \quad a : X^{\text{op}} \otimes X \rightarrow V.$$

To the mate $y = \lceil a \rceil : X \rightarrow (|X| \dashv\dashv V)$ of the first \mathcal{T} -functor we refer as the Yoneda functor. We have the following

Theorem 3.9 ([CH07]). *Let $X = (X, a)$ be a \mathcal{T} -category. Then the following assertions hold.*

- (1) For all $\mathfrak{x} \in TX$ and $\varphi \in (|X| \dashv\dashv V)$, $\llbracket m_X, \text{hom}_{\xi} \rrbracket (Ty(\mathfrak{x}), \varphi) \leq \varphi(\mathfrak{x})$.
- (2) Let $\varphi \in (|X| \dashv\dashv V)$. Then

$$\forall \mathfrak{x} \in TX. \varphi(\mathfrak{x}) \leq \llbracket m_X, \text{hom}_{\xi} \rrbracket (Ty(\mathfrak{x}), \varphi) \iff \varphi : X^{\text{op}} \rightarrow V \text{ is a } \mathcal{T}\text{-functor}.$$

Consequently, we put $\hat{X} = (\hat{X}, \hat{a})$ where

$$\hat{X} = \{\psi \in (|X| \dashv\dashv V) \mid \psi : X^{\text{op}} \rightarrow V \text{ is a } \mathcal{T}\text{-functor}\}$$

considered as a subcategory of $|X| \dashv\dashv V$. In particular, $y : X \rightarrow \hat{X}$ is full and faithful.

Example 3.10. For $X \in \mathcal{U}_2\text{-Cat} \cong \text{Top}$, $\psi \in \hat{X}$ is the characteristic function of a Zariski closed and down-closed subset $\mathcal{A} \subseteq UX$ (see Examples 3.4). We will give now an alternative description of \hat{X} , as the set $F_0(X)$ of (possibly improper) filters on the lattice τ of open sets of X , in terms of the bijective maps

$$\hat{X} \xrightarrow{\Phi} F_0(X) \quad \text{and} \quad F_0(X) \xrightarrow{\Pi} \hat{X},$$

where $\Phi(\mathcal{A}) = \bigcap \mathcal{A} \cap \tau$ and $\Pi(\mathfrak{f}) = \{x \in UX \mid \mathfrak{f} \subseteq x\}$. Clearly, $\mathcal{A} = \Pi(\mathfrak{f})$ is Zariski closed. If $x \leq \eta$ for some $x \in UX$ and $\eta \in \mathcal{A}$, then, for each $A \in x$ and $B \in \mathfrak{f}$, we have

$$\overline{A} \cap B \neq \emptyset$$

which, since B is open, gives $A \cap B \neq \emptyset$. Hence $\mathfrak{f} \subseteq x$, that is, $x \in \mathcal{A}$. Furthermore, one easily sees that $\mathfrak{f} = \Phi\Pi(\mathfrak{f})$ and $\mathcal{A} \subseteq \Pi\Phi(\mathcal{A})$. On the other hand, for $x \supseteq \bigcap \mathcal{A} \cap \tau$ and $A \in x$ we have $X \setminus \overline{A} \notin \bigcap \mathcal{A}$, and therefore $X \setminus \overline{A} \notin x$ for some $x \in \mathcal{A}$, hence $\overline{A} \in x$. Consequently, $\overline{A} \subseteq \bigcup \mathcal{A}$ and, since \mathcal{A} is Zariski closed, $x \leq \eta$ for some $\eta \in \mathcal{A}$. But \mathcal{A} is also down-closed, hence $x \in \mathcal{A}$. In a similar way (in fact, even easier) one can show that there are bijective maps

$$\check{X} \xrightarrow{\Phi'} F_1(X) \quad \text{and} \quad F_1(X) \xrightarrow{\Pi'} \check{X},$$

where $\check{X} = \{\mathcal{A} \subseteq UX \mid \mathcal{A} \text{ is Zariski closed and up-closed}\}$, $F_1(X)$ is the set of all (possibly improper) filters on the lattice σ of closed sets of X , $\Phi'(\mathcal{A}) = \bigcap \mathcal{A} \cap \sigma$ and $\Pi'(\mathfrak{f}) = \{x \in UX \mid \mathfrak{f} \subseteq x\}$. Furthermore, for any $\mathcal{A} \subseteq UX$ Zariski-closed, its down-closure $\downarrow\mathcal{A}$ is Zariski-closed as well. To see this, let $x \in \text{cl}(\downarrow\mathcal{A})$. Hence $x \in \bigcup \downarrow\mathcal{A}$ and therefore, for any $A \in x$, we have $\overline{A} \in \bigcup \mathcal{A}$. Define

$$\mathfrak{j} = \{B \subseteq X \mid \forall \alpha \in \mathcal{A}. B \notin \alpha\}.$$

Then \mathfrak{j} is an ideal, and $\mathfrak{j} \cap \{\overline{A} \mid A \in x\} = \emptyset$. Hence there is some $\eta \in UX$ such that $x \leq \eta$ and $\mathfrak{j} \cap \eta = \emptyset$. But the latter fact gives us $\eta \subseteq \bigcup \mathcal{A}$, that is, $\eta \in \text{cl} \mathcal{A} = \mathcal{A}$. We conclude $x \in \downarrow\mathcal{A}$. With a similar proof one can show that $\uparrow\mathcal{A}$ is Zariski closed for each Zariski closed subset $\mathcal{A} \subseteq UX$ (but now use $x \in \text{cl}(\uparrow\mathcal{A}) \iff \bigcap \uparrow\mathcal{A} \subseteq x$).

The topology in \hat{X} is the *compact-open topology*, which has as basic open sets

$$B(\mathcal{B}, \{0\}) = \{\mathcal{A} \in \hat{X} \mid \mathcal{A} \cap \mathcal{B} = \emptyset\}, \quad \mathcal{B} \subseteq UX \text{ Zariski closed.}$$

Since $B(\mathcal{B}, \{0\}) = B(\uparrow\mathcal{B}, \{0\})$, it is enough to consider Zariski closed and up-closed subsets $\mathcal{B} \subseteq UX$. Hence, using the bijections $\hat{X} \cong F_0(X)$ and $\check{X} \cong F_1(X)$, $F_0(X)$ has

$$\{\mathfrak{f} \in F_0(X) \mid \exists A \in \mathfrak{f}, B \in \mathfrak{g}. A \cap B = \emptyset\} \quad (\mathfrak{g} \in F_1(X))$$

as basic open sets. Clearly, it is enough to consider $\mathfrak{g} = \dot{B}$ the principal filter induced by a closed set B , so that all sets

$$\{\mathfrak{f} \in F_0(X) \mid \exists A \in \mathfrak{f}. A \cap B = \emptyset\} = \{\mathfrak{f} \in F_0(X) \mid X \setminus B \in \mathfrak{f}\} \quad (B \subseteq X \text{ closed})$$

form a basis for the topology on $F_0(X)$. But this is precisely the topology on $F_0(X)$ considered in [Esc97].

3.6. L-separation. We call a \mathcal{T} -category $X = (X, a)$ *L-separated* whenever the ordered set $\mathcal{T}\text{-Cat}(Y, X)$ is anti-symmetric, for each \mathcal{T} -category Y . The full subcategory of $\mathcal{T}\text{-Cat}$ consisting of all L-separated \mathcal{T} -categories is denoted by $\mathcal{T}\text{-Cat}_{\text{sep}}$.

Proposition 3.11. *Let $X = (X, a)$ be a \mathcal{T} -category. Then the following assertions are equivalent.*

- (i) X is L-separated.
- (ii) $x \cong y$ implies $x = y$, for all $x, y \in X$.
- (iii) For all $x, y \in X$, if $a(e_X(x), y) \geq k$ and $a(e_X(y), x) \geq k$, then $x = y$.
- (iv) $y : X \rightarrow \hat{X}$ is injective.

Proof. As for Proposition 1.5. □

Corollary 3.12. (1) *The \mathcal{T} -category $\mathbb{V} = (\mathbb{V}, \text{hom}_\varepsilon)$ is separated.*

- (2) *For all \mathcal{T} -categories $Y = (Y, b)$ and $X = (X, a)$ where Y is L-separated and $a \cdot T_\varepsilon a = a \cdot m_X$, Y^X is L-separated. In particular, $|X| \multimap \mathbb{V}$ is L-separated, for each \mathcal{T} -category X .*

(3) Any subcategory of a L -separated \mathcal{T} -category is L -separated. In particular, \hat{X} is L -separated, for each \mathcal{T} -category X .

Examples 3.13. A topological space is L -separated if and only if it is T_0 , whereas an approach space $X = (X, d)$ with distance function $d : PX \times X \rightarrow \mathbb{P}_+$ is L -separated if and only if

$$d(\{x\}, y) = 0 = d(\{y\}, x) \Rightarrow x = y$$

for all $x, y \in X$.

3.7. L-completeness. As in 1.7, we call a \mathcal{T} -category $X = (X, a)$ L -complete if every adjunction $\varphi \dashv \psi$ with $\varphi : Z \dashv X$ and $\psi : X \dashv Z$ is of the form $f_* \dashv f^*$ for a \mathcal{T} -functor $f : Z \rightarrow X$. Of course, f is up to equivalence uniquely determined by $\varphi \dashv \psi$, and is indeed unique if X is L -separated. As before, it is enough to consider $Z = E$ (see also [CH07])

Proposition 3.14. Let $X = (X, a)$ be a \mathcal{T} -category. The following assertions are equivalent.

- (i) X is L -complete.
- (ii) Each left adjoint \mathcal{T} -module $\varphi : E \dashv X$ is of the form $\varphi = x_*$ for some x in X .
- (iii) Each right adjoint \mathcal{T} -module $\psi : X \dashv E$ is of the form $\psi = x^*$ for some x in X .

A topological space is L -complete precisely if it is weakly-sober, that is, if each irreducible closed set is the closure of a point. A similar result holds for approach spaces: L -completeness is equivalent to each irreducible closed variable set A is representable (see [CH07] for details). Furthermore, in both cases we have that L -complete and L -separated (approach respectively topological) spaces are precisely the fixed objects of the dual adjunction with the category of (approach) frames induced by $\mathbb{V} = \mathbf{2}$ respectively $\mathbb{V} = \mathbb{P}_+$ (see [VO05] for details about the approach case).

For a pair $\psi : X \dashv Y$ and $\varphi : Y \dashv X$ of adjoint \mathcal{T} -modules $\varphi \dashv \psi$, the same calculation as in 1.7 shows that $\varphi = 1_X^* \circ \psi$. Since for each \mathcal{T} -module $\psi : X \dashv Y$ we have $(1_X^* \circ \psi) \circ \psi \leq 1_X^*$, ψ is right adjoint if and only if $\psi \circ (1_X^* \circ \psi) \geq (1_Y)_*$. Considering in particular $Y = E$, a \mathcal{T} -module $\psi : X \dashv E$ is right adjoint if and only if

$$k \leq \bigvee_{\mathfrak{x} \in TX} \psi(\mathfrak{x}) \otimes \xi \cdot T\varphi(\mathfrak{x})$$

where $\varphi = 1_X^* \circ \psi$. Note that $\bigvee \{\xi \cdot T\psi(\mathfrak{x}) \mid \mathfrak{x} \in TTX, m_X(\mathfrak{x}) = \mathfrak{x}\} = \psi(\mathfrak{x})$ since $\psi : |X| \rightarrow \mathbb{V}$ is a \mathcal{T} -functor, hence, with the help of Lemma 3.5, we see that

$$\begin{aligned} \varphi(x) &= \bigwedge_{\mathfrak{x} \in TX} \left(\left(\bigvee_{\substack{\mathfrak{x} \in TTX, \\ m_X(\mathfrak{x}) = \mathfrak{x}}} \xi \cdot T\psi(\mathfrak{x}) \right) \dashv a(\mathfrak{x}, x) \right) \\ &= \bigwedge_{\mathfrak{x} \in TX} (\psi(\mathfrak{x}) \dashv a(\mathfrak{x}, x)) \\ &= \hat{a}(e_{\hat{X}}(\psi), y(x)). \end{aligned}$$

Lemma 3.15. Let $\psi : X \dashv E$ be a \mathcal{T} -module and put $\varphi = 1_X^* \circ \psi$. Then, for each $\mathfrak{x} \in TX$,

$$\xi \cdot T\varphi(\mathfrak{x}) = T_\xi \hat{a}(e_{T\hat{X}} \cdot e_{\hat{X}}(\psi), T y(\mathfrak{x})).$$

Proof. Since $\xi \cdot T\varphi(\mathfrak{x}) = T_\xi \varphi(\mathfrak{x})$, the result follows from applying T_ξ to the equality above. \square

Hence we have

Proposition 3.16. *Let $X = (X, a)$ be \mathcal{T} -category. A \mathcal{T} -module $\psi : X \dashv\vdash E$ is right adjoint if and only if*

$$(\dagger) \quad k \leq \bigvee_{x \in TX} \psi(x) \otimes T_{\xi} \hat{a}(e_{T\hat{X}} \cdot e_{\hat{X}}(\psi), T y(x)).$$

Given a \mathcal{T} -category $X = (X, a)$, we call a \mathcal{T} -functor $\psi : |X| \rightarrow \mathbf{V}$ *tight* if $\psi : X^{\text{op}} \rightarrow \mathbf{V}$ is a \mathcal{T} -functor and if, considered as a \mathcal{T} -module $\psi : X \dashv\vdash E$, it is right adjoint, that is, if it satisfies (\dagger) .

Example 3.17. For a topological space X and $\psi \in \hat{X}$, as before we can identify ψ with a Zariski closed and down-closed subset $\mathcal{A} \subseteq UX$, and then $1_X^* \circ \psi$ with

$$A = \{x \in X \mid \forall a \in \mathcal{A}. a \rightarrow x\}.$$

Then ψ is tight if, and only if, there exists some $a \in \mathcal{A}$ with $A \in a$. Furthermore, under the bijection $\hat{X} \cong F_0(X)$ (see Example 3.10), a tight map ψ corresponds to a filters $\mathfrak{f} \in F_0(X)$ with $(\text{Lim } \mathfrak{f}) \# \mathfrak{f}$, where $\text{Lim } \mathfrak{f}$ denotes the set of all limit points of \mathfrak{f} and $A \# \mathfrak{g}$ if $\forall B \in \mathfrak{f}. A \cap B \neq \emptyset$. Furthermore, for each $\mathfrak{f} \in F_0(X)$ we have

$$(\text{Lim } \mathfrak{f}) \# \mathfrak{f} \iff \mathfrak{f} \text{ is completely prime,}$$

that is, if $\bigcup_{i \in I} U_i \in \mathfrak{f}$, then $U_i \in \mathfrak{f}$ for some $i \in I$. In fact, if $(\text{Lim } \mathfrak{f}) \# \mathfrak{f}$ and $\bigcup_{i \in I} U_i \in \mathfrak{f}$ for some family of open subsets of X , then $(\text{Lim } \mathfrak{f}) \cap \bigcup_{i \in I} U_i \neq \emptyset$, and therefore, for some $i \in I$, U_i contains a limit point of \mathfrak{f} . Hence $U_i \in \mathfrak{f}$. On the other hand, assume that \mathfrak{f} is completely prime. Suppose that $U \in \mathfrak{f}$ does not contain a limit point of \mathfrak{f} . Then, for each $x \in U$, there is an open neighborhood U_x of x with $U_x \notin \mathfrak{f}$. But $\bigcup_{x \in X} U_x \in \mathfrak{f}$ and, since \mathfrak{f} is completely prime, $U_x \in \mathfrak{f}$ for some $x \in U$, a contradiction.

3.8. L-injectivity. The notions of L-dense \mathcal{T} -functor, L-equivalence as well as L-injective \mathcal{T} -category can now be introduced as in 1.8. More precise, we call a \mathcal{T} -functor $f : (X, a) \rightarrow (Y, b)$ *L-dense* if $f_* \circ f^* = 1_{\hat{X}}$, which amounts to $b \cdot T f \cdot T f^\circ \cdot T_{\xi} b \cdot m_Y^\circ = b$. L-dense \mathcal{T} -functors have the same composition-cancellation properties as \mathbf{V} -functors (see 1.8). A fully faithful L-dense \mathcal{T} -functor is an *L-equivalence*, which can be equivalently expressed by saying that f_* is an isomorphism in $\mathcal{T}\text{-Mod}$. A \mathcal{T} -category Z is called *pseudo-injective* if, for all \mathcal{T} -functors $f : X \rightarrow Z$ and fully faithful \mathcal{T} -functors $i : X \rightarrow Y$, there exists a \mathcal{T} -functor $g : Y \rightarrow Z$ such that $g \cdot i \cong f$. Z is called *L-injective* if this extension property is only required along L-equivalences $i : X \rightarrow Y$. Of course, for a L-separated \mathcal{T} -category Z , $g \cdot i \cong f$ implies $g \cdot i = f$, and then pseudo-injectivity coincides with the usual notion of injectivity. The following two results can be proven as in 1.8

Lemma 3.18. *The \mathcal{T} -category \mathbf{V} is injective.*

Proposition 3.19. *For all \mathcal{T} -categories $Y = (Y, b)$ and $X = (X, a)$ where Y is L-injective (pseudo-injective) and $a \cdot T_{\xi} a = a \cdot m_X$, Y^X is L-injective (pseudo-injective).*

In particular, we obtain the injectivity of the \mathcal{T} -category $|X| \dashv\vdash \mathbf{V}$. Later on we will see that also \hat{X} and \tilde{X} are L-injective.

4. L-CLOSURE

4.1. L-dense \mathcal{T} -functors. As in 2.1, L-dense \mathcal{T} -functors can be characterized as “epimorphisms up to \cong ”. However, we will use here a slightly different proof.

Lemma 4.1. *Let $X = (X, a)$ be a \mathcal{T} -category, $M \subseteq X$ and $i : M \hookrightarrow X$ the embedding of M into X . Then i is dense if and only if*

$$(\ddagger) \quad k \leq \bigvee_{a \in TM} a(a, x) \otimes T_{\xi} a(T e_X \cdot e_X(x), a)$$

for all $x \in X$.

Proof. Recall that i is dense whenever $i_* \circ i^* \geq a$, that is,

$$a(\mathfrak{x}, x) \leq \bigvee_{\alpha \in TM} \bigvee_{\substack{\mathfrak{x} \in TTX \\ m_X(\mathfrak{x})=x}} a(\alpha, x) \otimes T_\xi a(\mathfrak{x}, \alpha)$$

for all $\mathfrak{x} \in T X$ and $x \in X$. If i is dense, then (‡) follows from the inequality above by putting $\mathfrak{x} = e_X(x)$ and using $m_X^\circ \cdot e_X = e_{TX} \cdot e_X$ (see Subsection 3.1). On the other hand, from (‡) we obtain

$$\begin{aligned} a(\mathfrak{x}, x) &\leq \bigvee_{\alpha \in TM} a(\alpha, x) \otimes T_\xi a(Te_X \cdot e_X(x), \alpha) \otimes a(\mathfrak{x}, x) \\ &\leq \bigvee_{\alpha \in TM} a(\alpha, x) \otimes T_\xi T_\xi a(e_{TTX} \cdot e_{TX}(\mathfrak{x}), e_{TX} \cdot e_X(x)) \otimes T_\xi a(Te_X \cdot e_X(x), \alpha) \\ &\leq \bigvee_{\alpha \in TM} a(\alpha, x) \otimes T_\xi a(e_{TX}(\mathfrak{x}), \alpha) \\ &\leq \bigvee_{\alpha \in TM} \bigvee_{\substack{\mathfrak{x} \in TTX \\ m_X(\mathfrak{x})=x}} a(\alpha, x) \otimes T_\xi a(\mathfrak{x}, \alpha). \end{aligned} \quad \square$$

Proposition 4.2. *For a \mathcal{T} -functor $i : M \rightarrow X$, the following assertions are equivalent.*

- (i) $i : M \rightarrow X$ is L -dense.
- (ii) For all \mathcal{T} -functors $f, g : X \rightarrow Y$, with $f \cdot i = g \cdot i$ one has $f \cong g$.
- (iii) For all \mathcal{T} -functors $f, g : X \rightarrow \mathbb{V}$, with $f \cdot i = g \cdot i$ one has $f = g$.

Proof. Assume first (i), i.e. $i : M \rightarrow X$ is L -dense. Then, from $f \cdot i = g \cdot i$ we obtain $f_* = g_*$ since $i_* \circ i^* = 1_X^*$. The implication (ii) \Rightarrow (iii) is trivially true. Assume now (iii). According to the remarks made above, we can assume that $i : M \rightarrow X$ is the embedding of a subset $M \subseteq X$. Let $x \in X$. First note that

$$\varphi : X \rightarrow \mathbb{V}, \quad y \mapsto a(e_X(x), y)$$

\mathcal{T} -functor since $a : |X| \otimes X \rightarrow \mathbb{V}$ is so. Using the same argument as in [Hof07, Lemma 6.8], we see that also

$$\psi : X \rightarrow \mathbb{V}, \quad y \mapsto \bigvee_{\mathfrak{x} \in TM} T_\xi a(Te_X \cdot e_X(x), \mathfrak{x}) \otimes a(\mathfrak{x}, y)$$

is a \mathcal{T} -functor. Clearly, for each $y \in X$ we have $\psi(y) \leq \varphi(y)$. If $y \in M$, we can choose $\mathfrak{x} = e_X(y) \in TM$ and therefore, using $Te_X \cdot e_X = e_{TX} \cdot e_X$ and op-laxness of e , obtain $\varphi(y) \leq \psi(y)$. Hence $\varphi|_M = \psi|_M$, and from our assumption (iii) we deduce $k \leq \varphi(x) = \psi(x)$. \square

4.2. L-closure. For a \mathcal{T} -category $X = (X, a)$ and $M \subseteq X$, we define the L -closure of M in X by

$$\overline{M} = \{x \in X \mid \forall f, g : X \rightarrow Y. (f|_M = g|_M \Rightarrow f(x) \cong g(x))\}.$$

Hence \overline{M} is the largest subset N of X making the inclusion map $i : M \hookrightarrow N$ dense.

Proposition 4.3. *Let $X = (X, a)$ be a \mathcal{T} -category, $M \subseteq X$ and $x \in X$. Then the following assertions are equivalent.*

- (i) $x \in \overline{M}$.
- (ii) $k \leq \bigvee_{\mathfrak{x} \in TM} a(\mathfrak{x}, x) \otimes T_\xi a(Te_X \cdot e_X(x), \mathfrak{x})$
- (iii) $i^* \circ x_* \dashv x^* \circ i_*$, where $i : M \hookrightarrow X$ is the inclusion map.
- (iv) $1_E^* \leq x^* \circ i_* \circ i^* \circ x_*$,
- (v) $i^* \circ x_* \dashv x^* \circ i_*$.
- (vi) $x_* : E \dashv \rightarrow X$ factors through $i_* : M \dashv \rightarrow X$ by a map $\varphi : E \dashv \rightarrow M$ in $\mathcal{T}\text{-Mod}$.

Proof. As for Proposition 2.2, using now Lemma 4.1. \square

We can now proceed as in 2.2.

Proposition 4.4. *For a \mathcal{T} -functor $f : X \rightarrow Y$ and $M, M' \subseteq X$, $N \subseteq Y$, we have*

- (1) $M \subseteq \overline{M}$ and $M \subseteq M'$ implies $\overline{M} \subseteq \overline{M'}$.
- (2) $\overline{\overline{M}} = \overline{M}$ and, if $T\emptyset = \emptyset$, then $\overline{\emptyset} = \emptyset$.
- (3) $f(\overline{M}) \subseteq \overline{f(M)}$ and $f^{-1}(\overline{N}) \supseteq \overline{f^{-1}(N)}$.
- (4) If k is \vee -irreducible and T preserves binary sums, then $\overline{M \cup M'} = \overline{M} \cup \overline{M'}$.

Corollary 4.5. *If k is \vee -irreducible in \mathbb{V} and T preserves finite sums, then the L-closure operator defines a topology on X such that every \mathcal{T} -functor becomes continuous. Hence, L-closure defines a functor $L : \mathcal{T}\text{-Cat} \rightarrow \text{Top}$.*

Example 4.6. For a topological space X , $x \in X$ lies in the L-closure of $A \subseteq X$ precisely if there exists some ultrafilter $\mathfrak{x} \in UA$ with $\bar{x} \in \mathfrak{x}$ and which converges to x ; in other words, for each neighborhood U of x we have $U \cap \bar{x} \cap A \neq \emptyset$. Hence the L-closure of a topological space X coincides with the so called *b-closure* [Bar68].

4.3. L-separatedness via the L-closure.

Proposition 4.7. *Let $X = (X, a)$ be a \mathcal{T} -category and $\Delta \subseteq X \times X$ the diagonal. Then*

$$\overline{\Delta} = \{(x, y) \in X \times Y \mid x \cong y\}$$

Proof. As for Proposition 2.6. □

Corollary 4.8. *A \mathcal{T} -category X is L-separated if and only if the diagonal Δ is closed in $X \times X$.*

Theorem 4.9. *$\mathcal{T}\text{-Cat}_{\text{sep}}$ is an epi-reflective subcategory of $\mathcal{T}\text{-Cat}$, where the reflection map is given by $y_X : X \rightarrow y_X(X)$, for each \mathcal{T} -category X . Hence, limits of L-separated \mathcal{T} -categories are formed in $\mathcal{T}\text{-Cat}$, while colimits are obtained by reflecting the colimit formed in $\mathcal{T}\text{-Cat}$. The epimorphisms in $\mathcal{T}\text{-Cat}$ are precisely the L-dense \mathcal{T} -functors.*

4.4. L-completeness via the L-closure.

Lemma 4.10. *Let $X = (X, a)$ be a \mathcal{T} -category and $M \subseteq X$. Then the following assertions hold.*

- (1) *Assume that X is L-complete and M be L-closed. Then M is L-complete.*
- (2) *Assume that X is L-separated and M is L-complete. Then M is L-closed.*

Proof. As for Lemma 2.9. □

Theorem 4.11. *Let $X = (X, b)$ be a \mathcal{T} -category. The following assertions are equivalent.*

- (i) *X is L-complete.*
- (ii) *X is L-injective.*
- (iii) *$y : X \rightarrow \tilde{X}$ has a pseudo left-inverse \mathcal{T} -functor $R : \tilde{X} \rightarrow X$, i.e. $R \cdot y \cong 1_X$.*

Proof. To see (i) \Rightarrow (ii), let $i : A \rightarrow B$ be a fully faithful dense \mathcal{T} -functor and $f : A \rightarrow X$ be a \mathcal{T} -functor. Since $i_* \dashv i^*$ is actually an equivalence of \mathcal{T} -modules, we have $f_* \circ i^* \dashv i_* \circ f^*$. Hence, since X is L-complete, there is a \mathcal{T} -functor $g : B \rightarrow X$ such that $g_* = f_* \circ i^*$, hence $g_* \circ i_* = f_*$.

The implication (ii) \Rightarrow (iii) is surely true since $y : X \rightarrow \tilde{X}$ is dense and fully faithful.

Finally, to see (iii) \Rightarrow (i), let $R : \tilde{X} \rightarrow X$ be a left inverse of $y : X \rightarrow \tilde{X}$. Then $y \cdot R = 1_{\tilde{X}}$ since $y : X \rightarrow \tilde{X}$ is dense and \tilde{X} is L-separated. Hence, for each right adjoint \mathcal{T} -module $\psi : X \rightarrow E$, we have $\psi = R(\psi)^*$. □

Therefore we have that $|X| \multimap V$ is L-complete. Our next result shows that also \hat{X} is L-complete.

Proposition 4.12. \hat{X} is L-closed in $|X| \multimap V$, for each \mathcal{T} -category X .

Proof. Let $X = (X, a)$ be a \mathcal{T} -category and assume that $\varphi \in (|X| \multimap V)$ belongs to the closure of \hat{X} , that is,

$$k \leq \bigvee_{u \in T\hat{X}} \llbracket m_X, \text{hom}_\xi \rrbracket(u, \varphi) \otimes T_\xi \llbracket m_X, \text{hom}_\xi \rrbracket(Te_{|X| \multimap V} \cdot e_{|X| \multimap V}(\varphi), u).$$

We wish to show that $r(\mathfrak{x}, \mathfrak{y}) \otimes \varphi(\mathfrak{y}) \leq \varphi(\mathfrak{x})$ for all $\mathfrak{x}, \mathfrak{y} \in TX$, where $r = T_\xi a \cdot m_X^\circ$.

First note that, for all $\alpha, \beta \in (|X| \multimap V)$,

$$e_{|X| \multimap V}^\circ \cdot \llbracket m_X, \text{hom}_\xi \rrbracket(\alpha, \beta) = \bigwedge_{\mathfrak{x} \in TX} (\alpha(\mathfrak{x}) \multimap \beta(\mathfrak{x})).$$

Hence, with $h_{\mathfrak{x}} : (|X| \multimap V) \rightarrow (|X| \multimap V)$, $h_{\mathfrak{x}}(\alpha, \beta) = (\alpha(\mathfrak{x}) \multimap \beta(\mathfrak{x}))$, we have $Te_{|X| \multimap V}^\circ \cdot T_\xi \llbracket m_X, \text{hom}_\xi \rrbracket \leq T_\xi h_{\mathfrak{x}}$. Since the diagram

$$\begin{array}{ccc} (|X| \multimap V) \times (|X| \multimap V) & \xrightarrow{\text{ev}_x \times \text{ev}_x} & V \times V \\ & \searrow h_x & \downarrow \multimap \\ & & V \end{array}$$

commutes and in

$$\begin{array}{ccccc} T((|X| \multimap V) \times (|X| \multimap V)) & \xrightarrow{T(\text{ev}_x \times \text{ev}_x)} & T(V \times V) & \xrightarrow{T(\multimap)} & TV \\ \text{can} \downarrow & & \text{can} \downarrow & & \downarrow \xi \\ T(|X| \multimap V) \times T(|X| \multimap V) & \xrightarrow{T \text{ev}_x \times T \text{ev}_x} & TV \times TV \geq & & \downarrow \xi \\ & & \xi \times \xi \downarrow & & \downarrow \xi \\ & & V \times V & \xrightarrow{\multimap} & V \end{array}$$

the left hand side diagram commutes and in the right hand side diagram we have “lower path” greater or equal “upper path”, we have

$$T_\xi h_{\mathfrak{x}}(u, v) \leq (u(\mathfrak{x}) \multimap v(\mathfrak{x}))$$

for each $\mathfrak{x} \in TX$ and $u, v \in T(|X| \multimap V)$, where $u(\mathfrak{x}) = \xi \cdot T \text{ev}_x(u)$. Accordingly, $e_{|X| \multimap V}(\varphi)(\mathfrak{x}) = \varphi(\mathfrak{x})$ and we obtain

$$\forall \mathfrak{x} \in TX . T_\xi \llbracket m_X, \text{hom}_\xi \rrbracket(Te_{|X| \multimap V} \cdot e_{|X| \multimap V}(\varphi), u) \leq (\varphi(\mathfrak{x}) \multimap u(\mathfrak{x})).$$

Furthermore, for all $\mathfrak{x}, \mathfrak{y} \in TX$ we have

$$\begin{array}{ccc} \hat{X} & \xrightarrow{\Delta} & \hat{X} \times \hat{X} \xrightarrow{\text{ev}_x \times \text{ev}_y} V \times V \\ ! \downarrow & & \downarrow \multimap \\ 1 & \xrightarrow{r(\mathfrak{x}, \mathfrak{y})} & V \end{array} \quad \leq$$

and we obtain

$$r(\mathfrak{x}, \mathfrak{y}) \leq \xi \cdot T(\multimap) \cdot T(\text{ev}_x \times \text{ev}_y) \cdot T\Delta(u) \leq (u(\mathfrak{y}) \multimap u(\mathfrak{x}))$$

for each $u \in T\hat{X}$. We conclude that

$$\begin{aligned}
r(\mathfrak{x}, \mathfrak{y}) \otimes \varphi(\mathfrak{y}) &\leq \bigvee_{u \in T\hat{X}} r(\mathfrak{x}, \mathfrak{y}) \otimes \varphi(\mathfrak{y}) \otimes \llbracket m_X, \text{hom}_\xi \rrbracket(u, \varphi) \otimes T_\xi \llbracket m_X, \text{hom}_\xi \rrbracket(Te_{|X| \rightarrow \mathcal{V}} \cdot e_{|X| \rightarrow \mathcal{V}}(\varphi), u) \\
&\leq \bigvee_{u \in T\hat{X}} r(\mathfrak{x}, \mathfrak{y}) \otimes \varphi(\mathfrak{y}) \otimes (\varphi(\mathfrak{y}) \multimap u(\mathfrak{y})) \otimes \llbracket m_X, \text{hom}_\xi \rrbracket(u, \varphi) \\
&\leq \bigvee_{u \in T\hat{X}} r(\mathfrak{x}, \mathfrak{y}) \otimes u(\mathfrak{y}) \otimes \llbracket m_X, \text{hom}_\xi \rrbracket(u, \varphi) \\
&\leq \bigvee_{u \in T\hat{X}} u(\mathfrak{x}) \otimes \llbracket m_X, \text{hom}_\xi \rrbracket(u, \varphi) \leq \bigvee_{u \in T\hat{X}} u(\mathfrak{x}) \otimes (u(\mathfrak{x}) \multimap \varphi(\mathfrak{x})) \leq \varphi(\mathfrak{x}). \quad \square
\end{aligned}$$

Proposition 4.13. *Let $X = (X, a)$ be a \mathcal{T} -category and $\psi \in \hat{X}$. Then ψ is a right adjoint \mathcal{T} -module if and only if $\psi \in \overline{y(X)}$.*

Proof. By Proposition 3.16 and Theorem 3.9, ψ is right adjoint if and only if

$$k \leq \bigvee_{\mathfrak{x} \in TX} \hat{a}(T y(\mathfrak{x}), \psi) \otimes T_\xi \hat{a}(T e_{\hat{X}} \cdot e_{\hat{X}}(\psi), T y(\mathfrak{x})),$$

which means precisely that $\psi \in \overline{y(X)}$. □

The proposition above identifies \hat{X} as the L-closure of $y(X)$ in \hat{X} , and therefore as an L-complete \mathcal{T} -category. Furthermore, $y : X \rightarrow \hat{X}$ is fully faithful and L-dense. Hence we have

Theorem 4.14. *The full subcategory $\mathcal{T}\text{-Cat}_{\text{cpl}}$ of $\mathcal{T}\text{-Cat}_{\text{sep}}$ of L-complete \mathcal{T} -categories is an epi-reflective subcategory of $\mathcal{T}\text{-Cat}_{\text{sep}}$. The reflection map of a L-separated \mathcal{T} -category X is given by any full L-dense embedding of X into an L-complete and L-separated \mathcal{T} -category, for instance by $y : X \rightarrow \hat{X}$.*

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